

ALGORITHMS FOR COMPUTER PROGRAMS

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pour Catherine

ABSTRACT. This paper collects the algorithms that I developed for calculating the ζ -functions of some particular curves. In particular, the TATE algorithm for handling WEIERSTRASS equations of elliptic curves is treated.

PREFACE

This is an old collection of routines for doing elementary arithmetic on a computer, including calculations in finite fields. Some have emerged from some `Algol` programs of mine written in 1978–1980 at the Faculty of Mathematics in Bielefeld. They are now implemented in `C` (interim versions in `REXX` have been removed).

These notes are meant to supplement the program logic, with emphasis put on the conceptual ideas behind those algorithms. The only change in version 1.7 is to the factorisation routine.

Berlin, December 15, 2012

Berndt E. Schwerdtfeger

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version 1.7.484, March 4, 2015.

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History.

- v0.1:** April 5, 1998. Initial draft document starting with the gcd-routine (*Euclidean* algorithm).
- v1.0:** December 30, 2002. First published version containing
- the gcd and extended gcd algorithms,
 - the factorisation into primes and prime table programs (based on sieve of ERATOSTHENES),
 - two initial JACOBI programs (in `Algo1`) that I developed in 1978 in support of my conjecture ([5]) about the structure of the JACOBI variety $J_0(64)$ of the modular curve $X_0(64) : x^4 + y^4 + z^4 = 0$ (for the proof see [6]).
 - the elliptic program code in C
- v1.1:** January 3, 2003. enhanced prime factorisation program, updated the elliptic code accordingly, enhanced the fiber type routine
- v1.2:** January 5, 2003. Added the remark that the sources of my programs are published under the GNU General Public License. Added also the remaining JACOBI programs.
- v1.3:** March 11, 2003. modified the elliptic program (display routine).
- v1.4:** April 5, 2003. Everything was put under CVS control. Updated the C-routines to latest level.
- Comme j'ai commencé cet article il y a cinq ans au jour exact, je l'ai dédié à *Catherine* pour son anniversaire.
- v1.5:** June 1, 2007. Added rcs package. Reference to BÉZOUT.
- v1.6:** December 3, 2010. Sources are now licensed under the Apache License, Version 2.0 (2005) <http://www.apache.org/licenses/LICENSE-2.0>.
- v1.7:** December 15, 2012. modified the factorisation routine to use `long long` integers.

1. GREATEST COMMON DIVISOR

The greatest common divisor d of the (positive, rational) integers a, b is defined by BÉZOUT's identity

$$a\mathbb{Z} + b\mathbb{Z} = d\mathbb{Z}$$

Of course, d is only unique up to the units in \mathbb{Z} , i.e. up to sign, but it is always understood that we take $d > 0$.

The *Euclidean* algorithm gives an effective way to modify the generators (a, b) by operations $g \in GL(2, \mathbb{Z})$. It rests on the integral division with remainder

$$a = q \cdot b + r \quad \text{with } 0 \leq r < b$$

Here q and r , subject to the above condition, are uniquely determined. In particular,

$$a\mathbb{Z} + b\mathbb{Z} = b\mathbb{Z} + r\mathbb{Z}$$

Let $g \in GL(2, \mathbb{Z})$ be the matrix $g = \begin{pmatrix} q & 1 \\ 1 & 0 \end{pmatrix}$. The *Euclidean* algorithm now reads $(a, b) = (b, r) \cdot g$ with $\det g = -1$.

We have finished our search, when b divides a ($r = 0$) and then the gcd is $d = b$. Otherwise ($r > 0$), we may start over again, now with the pair (b, r) . This process will come to an end, as we eventually must reach the case $r = 0$.

Let n be the number of steps and let now $g \in GL(2, \mathbb{Z})$ be the product of the above elementary matrix operations. We have

$$(a, b) = (d, 0) \cdot g$$

$$\det g = (-1)^n$$

This will be re-written for $g = \begin{pmatrix} u & v \\ w & x \end{pmatrix}$ as follows

$$a = u \cdot d$$

$$b = v \cdot d$$

$$a \cdot x - b \cdot w = (-1)^n d$$

Remark, by construction, that all matrix elements of g are ≥ 0 .

The iterative step in the program calculates the matrix product

$$\begin{pmatrix} q & 1 \\ 1 & 0 \end{pmatrix} \cdot \begin{pmatrix} u & v \\ w & x \end{pmatrix} = \begin{pmatrix} qu + w & qv + x \\ u & v \end{pmatrix}$$

1.1. GCD in REXX. The following REXX program implements this algorithm, calculating the matrix g .

```

/*
 *
 /* -----*
 //
 // gcd          greatest common divisor
 //
 //   Input: a, b   integers
 //   Output: gcd(a,b)
 //
 //
 // $Date: 2010-12-02 23:46:56 $
 // $Revision: 1.4 $
 //
 // -----
 //
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 //
 // -----*/

/* get the arguments */

Arg a b .

```

```

If a = '?' | a = '' then
do
  say ' '
  say 'Syntax is:  gcd a b '
  say ' '
  say '   Output is d = gcd(a,b), the factors of a and b'
  say '   and the linear relation d = ax + by'
  exit
end

main:

d = a; r = b
u = 1; v = 0; w = 0; x = 1; det = 1

do until r = 0
  q = d % r
  d0= r                      /* save old r   */
  r = d - q*r                /* new r     */
  d = d0                     /* new d     */
  u0= u; v0= v
  u = q*u + w; v = q*v + x  /* matrix    */
  w = u0      ; x = v0     /* ... product */
  det = - det
end

say ' '
say a '=' d '*' u
say b '=' d '*' v
say a '*' x '-' b '*' w '=' d*det

return

/*
*/

```

1.2. **Extended GCD in C.** The following C program implements the same algorithm, but using the inverse matrices $\begin{pmatrix} q & 1 \\ 1 & 0 \end{pmatrix}^{-1} = \begin{pmatrix} 0 & 1 \\ 1 & -q \end{pmatrix}$. Instead of the equation $(a, b) = (d, 0) \begin{pmatrix} u & v \\ w & x \end{pmatrix}$ one considers now $(u, v) \cdot \begin{pmatrix} u_1 & v_1 & t_1 \\ u_2 & v_2 & t_2 \end{pmatrix} = (u_0, v_0, t_0)$ in the notation of the C program below, with the ultimate goal that $v_0 = 0$ vanishes and that $u_0 = d$ is the $\gcd(u, v)$. See KNUTH, *The Art of Computer Programming* [1, vol. 2, §4.5.2], Algorithm X.

```

/*
*/
-----

Module:  xgcd.c

Description:

    Extended Euclidean Algorithm
    finding the greatest common divisor

```

Input: a, b integers

Output: $x*a + y*b = d$

\$Date: 2012-12-02 16:28:11 +0100 (Sun, 02 Dec 2012) \$
 \$Revision: 316 \$
 Version: 1.7

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 limitations under the License.

-----*/

```
#include <stdio.h>
#include <stdlib.h>

int main(void);

int main(void)
{
    long u, v, u0, v0, u1, v1, u2, v2, t0, t1, t2;
    long q;

    /* get the input */
    scanf("%d %d", &u, &v);

    u0 = u; v0 = v;      /* the following matrix is the inverse ../
    u1 = 1; v1 = 0;      /* .. of my rexx procedure gcd          */
    u2 = 0; v2 = 1;      /* (u,v) * mat(u1,v1 // u2,v2) = (u0,v0) */
                        /* u1*u + u2*v = u0 = gcd(u,v)          */
    while (v0) {         /* u = u0*|v2| , v = u0*|v1|          */
        q = u0 / v0;

        t0 = u0 - q*v0;
        t1 = u1 - q*v1;
        t2 = u2 - q*v2;

        u0 = v0;
        u1 = v1;
```

```

    u2 = v2;

    v0 = t0;
    v1 = t1;
    v2 = t2;

}

printf("gcd(%d,%d) = %d\n",u,v,u0);
printf("The coefficients in x*u + y*v = d are: x = %d, y = %d\n", u1, u2);
printf(".. the divisor are u/d = %d, v/d = %d\n", abs(v2), abs(v1));

return 0;
}

/*
*/

```

1.3. Factorization into primes. The following C program implements the algorithm A of KNUTH [1, vol. 2, §4.5.4]. I have made use of the sieve of ERATOSTHENES (see §2).

I make use of long long integers, i.e. the input parameter can be as large as $2^{64} - 1 = 18,446,744,073,709,551,615 > 18 \cdot 10^{18}$. For my practical purposes this is sufficient.

```

//
/* -----
Module: fact.c

Description:

    Factorisation into primes

    Input:  N integer to be factored (of type unsigned long long)
           maximal number handled is
            $2^{64}-1 = 18,446,744,073,709,551,615 > 18E18$ 

    Output: p[] array of primes, e[] array of exponents

$Date: 2012-12-02 16:28:11 +0100 (Sun, 02 Dec 2012) $
$Revision: 316 $
Version: 1.7 changed prime array to long long

-----

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```

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```

-----*/

#include <stdio.h>
#include <inttypes.h>

int main(void);
int main(void)

{
// some "magic constants"

#define L      14000
#define M      2*L*(L+1)          // = 1152048000 (1152MB)
// #define M      1316300
#define YES    1                  // flag for primes
#define NO     0                  // flag for composite numbers

unsigned long s,t;
// unsigned long long int N;          // integer to be factorized
// unsigned long long p[30];         // array of prime factors @ 1.7
// int e[30];                        // array of exponents
unsigned long d[220000000];         // array of test divisors (880MB)

int sieve(unsigned long []);
// int fact( unsigned long long, unsigned long [], unsigned long long [], int []);
//                                     // @ 1.7
t = sieve(d);                       // t = 192080, p = 2632601 > 2^21 = 2097152

// for (s=0;s<t;s++)
//   printf("%lu%c",d[s], s%8 ? '\t': '\n');

printf("\nt = %lu, p = %lu",t,d[t-1]);
printf("\nL = %lu, 2*L*(L+1)=%lu, M = %lu",L,2*L*(L+1),M);
return 0;                            // EARLY EXIT !!!
// *****
// scanf("%llu", &N);

// printf("\nN = ");

// if (N) {
//   t = fact(N, d, p, e);              // factorize the number
//   printf("%llu = ",N);
//   if (t) {
//     for (s=1; s<t ; s++)
//       printf("%llu(%u) * ",p[s], e[s]); // @ 1.7
//     printf("%llu(%u)",p[t], e[t]);     // @ 1.7
//   }
//   else
//     printf("1");                      // if no factors, it is 1

```

```

// } else
//   printf("0");

// return 0;
}

/* -----
**
** Subroutine  sieve
**
** Description: int sieve(unsigned long d[]);
**
**      Generating array of test primes via sieve of Eratosthenes
**
**      Input:  none
**
**      Output: return value, vector d[] of test divisors
** -----*/

int sieve(unsigned long d[])
{
// sieve of Eratosthenes

    unsigned int a, i;                // loop ind, exponent

    /* the array prime[x] test for primality of 2*x + 1 */

    char prime[M+1];                 // byte array for flags

    for (a = 0; a <= M; a++)          // initialize ...
        prime[a] = YES;              // ... to everything prime
    for (a = 1; a <= L; a++)          // go thru all primes and
        if (prime[a])                // throw out the multiples
            for (i = 2*a*(a+1); i <= M; i += 2*a+1)
                prime[i] = NO;        // multiples are not prime

// populate the vector of test divisors

    i = 0;
    d[i++] = 2;                       // set the even prime

    for (a = 1; a <= M; a++)
        if (prime[a])
            d[i++] = 2*a + 1;         // set the uneven primes

    return i;
}

/* -----
**
** Subroutine  fact
**
** Description:
**      The number N is factored into primes

```



```

**      N = p[1]**e[1] * ... * p[t]**e[t]
**
**      Input:  N (integer), test divisors in vector d[]
**
**      Output: vector p[], vector e[]
**              the vector p[] contains the different primes
**              the vector e[] contains the orders of N at these primes
**              p[0] = t = return value = number of primes (also t = 0)
**              (if N = 0 then e[0] = 1, otherwise e[0] = 0)
**
** -----*/
/*
int fact(unsigned long long N, unsigned long d[], unsigned long long p[], int e[])
{
//                                     // @ 1.7

// this is Knuth's algorithm, modified with exponents

    unsigned int i, t, k;
    unsigned long long n, q, r;

    n = N;
    i = t = k = 0;

    while (n > 1) {
// A2 n = 1 ?
        q = n / d[k];
// A3 divide
        r = n - q * d[k];
        if (r == 0) {
// A4 zero remainder ?
            if (i == 0)
                p[+t] = d[k];
// A5 new factor found
                e[t] = ++i;
// increase exponent
                n = q;
            }
// return to A2
        else {
            if (q > d[k]) {
// A6 low quotient ?
                k = k + 1;
// no, .. try next
                i = 0;
            }
// return to A3
            else {
// yes, ..
                p[+t] = n;
// A7 n is prime
                e[t] = 1;
// exponent is 1
                n = 1;
// terminate
            }
        }
    }

    p[0] = t;
// number of factors
    if (n)
        e[0] = 0;
// e[0] discriminates
    else
        e[0] = 1;
// .. if n vanishes

    return t;
} */

//

```

2. ERATOSTHENES

//

/* -----

Module: prim.c

Description:

Eratosthenes (274-194 b.C.) used the following reasoning for obtaining all primes:

```

go thru all natural numbers q > 1
'flag' all multiples (2q, 3q, 4q ...)
stop at any time (number N, say).
all numbers not flagged up to N are prime
(since not a multiple of any other number)

```

This simple reasoning can be modified in the following way:

First only odd numbers need to be considered ($q=2a+1$). So the flags for the a's make up an array of half the size `prim[M+1]`, where the odd number $N = 2*M + 1$. Second, each time we only need to flag multiples of q: $q*q$, $q*(q+2)$, $q*(q+4)$, ... up to N (q is odd), since smaller multiples have already been accounted for.

This also means that q only loops from 3 up to \sqrt{N} . In the formula below we assume \sqrt{N} is integral, odd, i.e. $N = 2*M + 1 = (2*L+1)*(2*L+1)$, therefore $M = 2*L*(L+1)$.

This program uses the maximum $L = 127$ for the current implementation (memory model). Either other memory models or usage of bitflags instead of byte arrays can generate larger tables. This is done elsewhere (see `primx.c`).

\$Date: 2013-01-06 11:12:37 \$

\$Revision: 1.5 \$

History:

```

06.08.1989  1.00  initial version
05.12.1989  1.01  use the "? : " construct for deciding newline
31.01.1993  1.01.01 provide a prolog and description
19.02.1994  1.02  modified algorithm
20.02.1994  1.02.01 fully document the algorithm in the description
30.03.1997  1.02.02 copyright statement
05.04.2003  1.1  put into CVS repository
06.01.2013  1.5  modified description, no program change

```

Notes: adopted from an original Algol60 program written in 1979

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```

-----*/
#define L      127          // has to be constant
#define M      2*L*(L+1)    // = 35512, N = 65025 = 255*255
#define COLUMNS 8         // for the screen output
#define ON     0            // flag for primes
#define OFF    1            // flag for composite numbers

#include <stdio.h>

int main(void)
{
    unsigned int  a,b;      // loop ind
    char prim[M+1];        // byte array for flags

    for (a = 0; a <= M; ++a) // initialize ...
        prim[a] = ON;       // ... to everything prime
    for (a = 1; a <= L; ++a) // go thru all primes and
        if ( prim[a] == ON ) // throw out the multiples
            for (b = 2*a*(a+1); b <= M; b += 2*a+1)
                prim[b] = OFF; // multiples are not prime

    b = 1;
    printf("%u%c",2,'\t'); // 2 is the first prime
    for (a = 1; a <= M; ++a) // put out the result
        if ( prim[a] == ON ) // if it is a prime
            printf("%u%c", 2*a + 1 , ++b % COLUMNS ?'\t':'\n');
    return (0);
}

//
//
/* -----*/

```

Module: primx.c

Description:

The general reasoning is explained in prim.c, which uses a byte array for the flags, whereas here I use a bit-array (put into an array of long's).

I use 32-bit flat address space. On my particular machine with 2 GByte of storage the largest prime generated was

p = 4,x??,???,??? (

\$Date: 2010-12-03 00:15:38 \$

\$Revision: 1.4 \$

History:

```

21.05.1979          original version in Algol60 written in Bielefeld
20.02.1994  1.00    rewrite in C
21.02.1994  1.00.01 loop restriction added (overflow fixed)
30.03.1997  1.00.02 copyright notice
28.07.2001  1.00.03 adaption to GNU/Debian platform
30.03.2003  1.01.04 M = 520128, max prime = 33,288,251
21.01.2013  1.5     M = 226-1 = 67108863, max prime =

```

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```

-----*/
#define M      67108863          // 226-1          @1.5
#define N      33554432          // 225
#define COLUMNS 256            // for file output

#include <stdio.h>

int main(void)
{
    unsigned long p, q, r, y, z;          // loop ind
    unsigned long L[M+1];                // long array for bitflags

    for (p = 0; q <= M; ++p)              // initialize ...
        L[p] = 0;                          // ... to everything prime

    for (q = 0; 32*q*q <= M; ++q)          // go thru all flag fields .. @1.00.01
        for (r = 0; r <= 31; ++r)         // .. for all bits in a long
            if (!(1 & L[q]>>r))             // p = 64*q+2*r+3 is prime
            {
                y = 2*q*(32*q+2*r+3) + (2*r*(r+3)+3)/32;
                z = (2*r*(r+3)+3)%32;
                while ( 32*y + z <= (long) 32*M + 31 )
                {
                    L[y] |= (long)1 << z;    // flag as composite
                    y += 2*q+(z+2*r+3)/32;
                    z = (z+2*r+3)%32;
                }
            }

    z = 1;
    printf("%u%c ", 2, ',', ',');          // 2 is the first prime

```

```

FILE *fp = fopen("first.txt","w");

for (q = 0; q <= N; ++q)           // go thru all flag fields ..
  for (r = 0; r <= 31; ++r)       // .. for all bits in a long
    if (!(1 & L[q]>>r))           // the r-bit in L[q] is off (prime!)
      fprintf(fp,"%lu%c ", 64*q+2*r+3, ++z % COLUMNS ?',':'\n');
fclose(fp);

FILE *zp = fopen("zweit.txt","w");
for (q = N; q <= M; ++q)         // go thru all flag fields ..
  for (r = 0; r <= 31; ++r)       // .. for all bits in a long
    if (!(1 & L[q]>>r))           // the r-bit in L[q] is off (prime!)
      fprintf(zp,"%lu%c ", 64*q+2*r+3, ++z % COLUMNS ?',':'\n');
fprintf(zp,"\nt = %lu",z);
fclose(zp);
return (0);
}

//

```

3. THE JACOBI PROGRAMS

The two curves, that I was concerned with, are

$$\begin{aligned} \mathbf{E}: & y^2 = x^3 + x \\ \mathbf{C}: & x^4 + y^4 + z^4 = 0 \end{aligned}$$

Consult my article ‘The Zeta–Function of $x^4 + y^4 + z^4 = 0$ ’ [6] for the theory and more background.

The JACOBI programs calculate the number of points of them over the prime field \mathbb{F}_p and the quadratic and cubic extensions \mathbb{F}_{p^2} , \mathbb{F}_{p^3} .

There are these Algol programs to do the job:

JAC1: calculations over \mathbb{F}_p
JAC2: calculations over \mathbb{F}_{p^2}
JAC2A: calculations over \mathbb{F}_{p^2} for primes $p \equiv 3 \pmod{4}$.
JAC3: calculations over \mathbb{F}_{p^3}

Although it is not mentioned in the sources, they are provided on the basis of the Apache License (the *leader dots* are added for better legibility to the \TeX version of the programs).

3.1. Jac1: calculations over \mathbb{F}_p .

```

'begin''comment' JAC1 ;
. 'integer' n,q,p,e1,c1,u1,s1;
. 'integer''procedure' mod(a);
. 'value' a;
. 'integer' a;
. 'begin'
. . a := a - entier(a/p)*p ;
. . 'if' a 'greater' (p-1)/2 'then'
. . . a := a-p ;
. . mod := a
. 'end'; 'comment'

```

Es folgen die Anweisungen;

```

. output(4,('4B,('P E1 C1)'),9B,('U1 S1'))');
. input(1,('B3D'),n); 'for' a := 1 'step' 1 'until' n 'do'
. 'begin'
. . input(1,('B4D'),p);
. . 'begin'
. . . 'integer''array' q[-(p-1)/2:(p-1)/2];
. . . 'integer' x,y;
. . . 'for' x := -(p-1)/2 'step' 1 'until' (p-1)/2 'do' q[x] := 0;
. . . 'for' x := 0 'step' 1 'until' (p-1)/2 'do'
. . . . 'begin'
. . . . . y := mod(x*x);
. . . . . q[y] := 1
. . . . . 'end';
. . . 'for' x := -(p-1)/2 'step' 1 'until' (p-1)/2 'do'
. . . . 'begin'
. . . . . y := mod(x*mod(x*x+1));
. . . . . 'if' q[y] 'equal' 1 'then' e1 := e1 + 2
. . . . . 'end';
. . . 'for' x := -(p-1)/2 'step' 1 'until' (p-1)/2 'do' q[x] := 0;
. . . 'for' x := 0 'step' 1 'until' (p-1)/2 'do'
. . . . 'begin'
. . . . . y := mod(x*x);
. . . . . y := mod(y*y);
. . . . . q[y] := 1
. . . . . 'end';
. . . 'for' x := -(p-1)/2 'step' 1 'until' (p-1)/2 'do'
. . . . 'begin'
. . . . . y := mod(x*x);
. . . . . y := -mod(mod(y*y)+1);
. . . . . 'if' q[y] 'equal' 1 'then' c1 := c1 + 4
. . . . . 'end';
. . . 'if' entier((p-1)/8)*8 + 1 'equal' p 'then' c1 := c1 - 8;
. . . u1 := 1 + p - e1;
. . . s1 := 1 + p - c1;
. . . output(4,('/4ZD,4ZD,4ZD,-9ZD,-3ZD'),p,e1,c1,u1,s1);
. . . 'if' s1 'notequal' 3*u1 'then'
. . . . output(4,('5B,('Gegenbeispiel'))');
. . . 'if' 6*sqrt(p) - abs(s1) 'less' 1 'then'
. . . . output(4,('4B,ZZD.DD'),6*sqrt(p))
. . 'end'
. 'end'
'end' Programmende von JAC1

```

3.2. Jac2: calculations over \mathbb{F}_{p^2} .

```

'begin''comment' JAC2 ;
. 'integer' p,d,e,c,n,l;
. 'integer''procedure' mod(a);
. 'value' a;
. 'integer' a;
. 'begin'
. . a := a - entier(a/p)*p ;
. . 'if' a 'greater' (p-1)/2 'then'
. . . a := a-p ;
. . . mod := a
. 'end'; 'comment'

```

```

Anweisungen;
. output(4,('*4B,('P')8B('E2')8B('C2')8B('U2')8B('S2')')));
. input(1,('BDD'),n); 'for' l := 1 'step' 1 'until' n 'do'
. 'begin'
. . 'integer' q;
. . input(1,('B3D'),p);
. . q := (p-1)/2;
. . 'begin'
. . . 'integer''array' a[-q:q,-q:q];
. . . 'integer' x,y,t,u,v;
. . . 'for' x := -q 'step' 1 'until' q 'do' a[x,0] := 0;
. . . 'for' x := 0 'step' 1 'until' q 'do'
. . . . 'begin'
. . . . . y := mod(x*x);
. . . . . a[y,0] := 1
. . . . . 'end';
. . . . e := -2;
. . . . 'for' x := -q 'step' 1 'until' q 'do'
. . . . . 'begin'
. . . . . . y := mod(x*mod(x*x+1));
. . . . . . 'if' a[y,0] 'equal' 1 'then' e := e + 2
. . . . . . 'end';
. . . . . e := e*((1+p)*2 - e);
. . . . . output(4,('/4ZD,9ZD'),p,e);
. . . . . e := 1 + p*p - e;
. . . . . 'for' x := -q 'step' 1 'until' q 'do'
. . . . . . 'if' a[x,0] 'equal' 0 'then''goto' nr;
. nr:.. d := x;
. . . . . 'for' x := -q 'step' 1 'until' q 'do'
. . . . . . 'for' y := -q 'step' 1 'until' q 'do' a[x,y] := 0;
. . . . . . 'for' x := -q 'step' 1 'until' q 'do'
. . . . . . . 'for' y := -q 'step' 1 'until' q 'do'
. . . . . . . . 'begin'
. . . . . . . . . t := mod(x*x + y*y*d);
. . . . . . . . . v := mod(2*x*y);
. . . . . . . . . u := mod(t*t + v*v*d);
. . . . . . . . . v := mod(2*t*v);
. . . . . . . . . a[u,v] := 1
. . . . . . . . . 'end';
. . . . . . . c := -8
. . . . . . . 'for' x := -q 'step' 1 'until' q 'do'
. . . . . . . . 'for' y := -q 'step' 1 'until' q 'do'
. . . . . . . . . 'begin'
. . . . . . . . . . t := mod(x*x + y*y*d);
. . . . . . . . . . v := mod(2*x*y);
. . . . . . . . . . u := - mod(t*t + v*v*d + 1);
. . . . . . . . . . v := - mod(2*t*v);
. . . . . . . . . . 'if' a[u,v] 'equal' 1 'then' c := c + 4
. . . . . . . . . . 'end';
. . . . . . . output(4,('9ZD,-8ZD'),c,e);
. . . . . . . c := 1 + p*p - c;
. . . . . . . output(4,('8ZD'),c);
. . . . . . . 'if' c 'notequal' 3*e 'then'
. . . . . . . . output(4,('('Gegenbeispiel'))');
. . . . . . . 'if' 6*p - abs(c) 'less' 1 'then'
. . . . . . . . output(4,('ZZD.DD'),6*p);
. . . . . . . 'end'

```

```
. 'end'
'end' Programmende von JAC2
```

3.3. Jac2a: calculations over \mathbb{F}_{p^2} for $p \equiv 3 \pmod{4}$.

```
'begin''comment' JAC2A fuer Primzahlen kongruent 3 modulo 4 ;
. 'integer' p, e, c, n, q ;
. 'integer''procedure' mod(a); 'value' a ; 'integer' a ;
. 'begin'
. . a := a - entier(a/p)*p ;
. . 'if' a 'greater' (p-1)/2 'then'
. . . a := a-p ;
. . . mod := a
. 'end'; 'comment'
Anweisungen; n := 0;
. output(4,('4B','P')8B('E2')8B('C2')8B('U2')8B('S2')));
marke:
. input(1,('5BD'),p); q := (p-1)/2;
. 'begin''integer''array' a[-q:q, -q:q] ;
. . 'integer' x,y,t,u,v;
. . e := (1+p)*(1+p) ; output(4,('4ZD,9ZD'),p,e);
. . e := -2*p;
. . 'for' x := -q 'step' 1 'until' q 'do'
. . . 'for' y := -q 'step' 1 'until' q 'do' a[x,y] := 0;
. . 'for' x := -q 'step' 1 'until' q 'do'
. . . 'for' y := -q 'step' 1 'until' q 'do'
. . . . 'begin'
. . . . . t := mod(x*x - y*y); v := mod(2*x*y);
. . . . . u := mod(t*t - v*v); v := mod(2*t*v);
. . . . . a[u,v] := 1
. . . . 'end';
. . c := -8
. . 'for' x := -q 'step' 1 'until' q 'do'
. . . 'for' y := -q 'step' 1 'until' q 'do'
. . . . 'begin'
. . . . . t := mod(x*x - y*y); v := mod(2*x*y);
. . . . . u := - mod(t*t - v*v + 1); v := - mod(2*t*v);
. . . . . 'if' a[u,v] 'equal' 1 'then' c := c + 4
. . . . 'end';
. . output(4,('9ZD,-8ZD'),c,e);
. . c := 1 + p*p - c;
. . output(4,('8ZD'),c);
. . 'if' c 'notequal' 3*e 'then'
. . . output(4,('Gegenbeispiel'));
. . 'if' 6*p - abs(c) 'less' 1 'then'
. . . output(4,('ZZD.DD'),6*p);
. . n := n + 1;
. . 'if' n 'less' 25 'then''goto' marke
. 'end'
'end' Programmende von JAC2A
```

3.4. Jac3: calculations over \mathbb{F}_{p^3} .

```
'begin''comment' JAC 3 ;
. 'integer' p, q, e, c, d, s, t, n;
. 'integer''procedure' mod(a); 'value' a ; 'integer' a ;
. 'begin' a := a - entier(a/p)*p ; 'if' a 'greater' (p-1)/2 'then'
. . a := a-p ; mod := a 'end'; '
```



```

. 'procedure' vier(x,y,z,u,v,w) ; 'integer' x,y,z,u,v,w ;
. . 'begin'
. . . s := mod(x*x + 2* y*z*d) ;
. . . t := mod(2*x*y + 2*y*z + z*z*d) ;
. . . w := mod(2*x*z + y*y + z*z) ;
. . . u := mod(s*s + 2*t*w*d) ;
. . . v := mod(2*s*t + 2*t*w + w*w*d) ;
. . . w := mod(2*s*w + t*t + w*w) ;
. . 'end' ;
. output(4,('4B','P')8B('E3')8B('C3')8B('U3')8B('S3')));
. n := 0;
marke:
. input(1,('5BD,-DDB,DDD'),p,e,d); q := (p-1)/2;
. e := 1 + p*p*p + 3*p*e - e*e*e; output(4,('4ZD,9ZD'),p,e);
. e := 1 + p*p*p - e;
. 'begin''integer''array' l[-q:q, -q:q, -q:q];
. . 'integer' u, v, w, x, y, z;
. . 'for' x := -q 'step' 1 'until' q 'do'
. . . 'for' y := -q 'step' 1 'until' q 'do'
. . . . 'for' z := -q 'step' 1 'until' q 'do'
. . . . . l[x,y,z] := 0;
. . 'for' x := -q 'step' 1 'until' q 'do'
. . . 'for' y := -q 'step' 1 'until' q 'do'
. . . . 'for' z := -q 'step' 1 'until' q 'do'
. . . . . 'begin' vier(x,y,z,u,v,w) ; l[u,v,w] := 1 'end' ;
. . 'if' entier(q/4)*4 'equal' q 'then' c := -8 'else' c := 0 ;
. . 'for' x := -q 'step' 1 'until' q 'do'
. . . 'for' y := -q 'step' 1 'until' q 'do'
. . . . 'for' z := -q 'step' 1 'until' q 'do'
. . . . . 'begin'
. . . . . . vier(x,y,z,u,v,w) ; u := -mod(u+1) ; v := -v ; w := -w ;
. . . . . . 'if' l[u,v,w] 'equal' 1 'then' c := c + 4
. . . . . . 'end';
. . output(4,('9ZD,-8ZD'),c,e); c := 1 + p*p*p - c;
. . output(4,('8ZD'),c);
. . 'if' c 'notequal' 3*e 'then'
. . . output(4,('Gegenbeispiel')));
. . n := n + 1;
. . 'if' n 'less' 4 'then''goto' marke
. . 'end'
'end' Programmende von JAC 3

```

4. THE 'ELLIPTIC' PROGRAM

There are two Algol programs

ELLALG: version 3.3.1980 (without reduction)
 ELLIP: version 11.3.1980 (with incomplete reduction)

each of which contains some more routines. There are also versions from the 6th and 10th of March as well, which describe the development from ELLALG towards ELLIP, which I did not finish then. I was obviously trying to add the necessary modifications to calculate the minimal WEIERSTRASS equation.

Here is first the high level structure of ELLALG:

4.1. ELLALG. The main program does the following

```

read number of Weierstrass equations e
for k = 1 to e
do
. read the coefficients
. calculate the Tate parms (b2,b4,b6,b8,c4,c6,d)
. write out the equation (drop zero's)
. if singular (d=0) then
. . if c4 = 0 then write 'cusp' (additive)
. . . else write 'node' (multiplicative)
. . endif
. else (regular, that is rational elliptic curve)
. . write discriminant
. . call tp(c4,d,b,t) /* prime factor decomposition */
. . write instable fibers (ord(p,d) > 0)
. . . semistable fiber at p for ord(p,c4)=0
. . . instable fiber at p, potentially good for ord(p,j)>=0
. . . instable fiber at p, potentially bad for ord(p,j) <0
. endif
end

```

This is the list of subroutines

- ord(p,m) valuation (exponent) of m at prime p
- max(m,n) maximum of m and n
- sig(z) sign of z
- out(d) formatting routine
- tp(m,n,a,t) prime factor decomposition of the pair (m,n) into a prime factors, exponents in the matrix t

4.2. ELLIP. The list of subroutines is the same as above, plus the *Euclidean* algorithm mod(a,n,q,r) and the reduction routine red(w,x), which, in fact, is *not called*.

- ord(p,m) valuation (exponent) of m at prime p
- mod(a,n,q,r) Euclidean algorithm (one step) $a = q \cdot n + r$
- red(w,x) reduction routine
- max(m,n) maximum of m and n
- sig(z) sign of z
- out(d) formatting routine
- tp(m,n,a,t) prime factor decomposition of the rational m/n into a prime factors, exponents in the matrix t

It appears that the main routine is not completed: the calculation of a minimal WEIERSTRASS equation was never done (the pieces that appear in the code are commented out).

Nevertheless, this is what I have (the 'dots' are added here for easier reading of the program nesting):

```

'comment' ellip
Version: 11.3.1980
Weierstrass-Gleichungen ueber Z
singulaer: Knoten / Spitze
elliptisch:

```

```

J - Invariante
Diskriminante D
Singulaere Faser ueber p <=> P-ORD(D) > 0
Typ * semistabil                P-ORD(J.D) = 0
    + instabil, potentiell gut   P-ORD(J) > 0
    +* instabil, potentiell schlecht P-ORD(J.D) > 0, P-ORD(J) < 0

Fuehrer-Exponent F(P) = 1 bei Typ *
                    F(P) >= 2 bei Typ +, +* (und = 2, wenn p > 3)

minimale Gleichung ? ;

'begin'
. 'integer' a1,a2,a3,a4,a6,b2,b4,b6,b8,c4,c6,d,e,b,v;
. 'integer''array' t[0:2,1:15], w,x[1:5];
.
. 'integer''procedure' ord(p,m);
. 'begin'
. . 'integer' e;
. . e := 0;
. w: 'if' m/'p*p = m 'then'
. . 'begin'
. . . e := e + 1;
. . . m := m / p;
. . . 'goto' w
. . 'end';
. . ord := e
. 'end';
.
. 'procedure' mod(a,n,q,r);
. . 'integer' a, n, q, r;
. 'begin'
. . q := entier(a/n);
. . r := a - n*q;
. . 'if' n < 2*r 'then'
. . . 'begin'
. . . . q := q + 1;
. . . . r := r - n;
. . . 'end';
. 'end';
.
. 'procedure' red(w,x);
. . 'integer''array' w, x;
. 'begin'
. . 'integer' r, s, t;
. .
. . mod(w[1],2,s,x[1]);
. . mod(w[3] + s*w[1] - s*s,3,r,x[3]);
. . mod(w[2] - r*w[1],2,t,x[2]);
. . x[4] := w[4] + s*w[2] - 2*r*w[3] + (t - r*s)*w[1] + 3*r*r - 2*s*t;
. . x[5] := w[5] - r*w[4] + r*r*w[3] - r*r*r + t*w[2] - r*t*w[1] - t*t;
. 'end';
.
.
. 'integer''procedure' max(m,n);
. 'begin'
. . max := 'if' m < n 'then' n 'else' m;

```

```

. 'end';
.
. 'procedure' sig(z);
. 'begin'
. . 'if' z < 0 'then' output(4,<<<<- >>>>)
. . . 'else' output(4,<<<<+ >>>>)
. 'end';
.
. 'procedure' out(d);
. 'begin'
. . 'switch' s := s1,s2,s3,s4,s5,s6,s7,s8,s9,s10,s11,s12;
. . 'goto' s[entier(ln(d)/ln(10)) + 1];
. . s1: output(4,<<d>>,d); 'goto' p;
. . s2: output(4,<<2d>>,d); 'goto' p;
. . s3: output(4,<<3d>>,d); 'goto' p;
. . s4: output(4,<<4d>>,d); 'goto' p;
. . s5: output(4,<<5d>>,d); 'goto' p;
. . s6: output(4,<<6d>>,d); 'goto' p;
. . s7: output(4,<<7d>>,d); 'goto' p;
. . s8: output(4,<<8d>>,d); 'goto' p;
. . s9: output(4,<<9d>>,d); 'goto' p;
. . s10: output(4,<<10d>>,d); 'goto' p;
. . s11: output(4,<<11d>>,d); 'goto' p;
. . s12: output(4,<<12d>>,d); 'goto' p;
. . p:
. 'end';
.
. 'procedure' tp(m,n,a,t);
. 'integer' a; 'integer''array' t;
.
. 'begin'
. . 'integer' q,r,s;
. .
. . s := max(1,entier((sqrt(max(m,n)) - 1)/2));
. . 'begin'
. . . 'integer''array' l[1:s];
. . .
. . . 'for' q := 1 'step' 1 'until' s 'do' l[q] := 1;
. . . 'for' q := 1 'step' 1 'until' s 'do'
. . . . 'if' l[q] = 1 'then'
. . . . . 'begin'
. . . . . . 'for' r := 3*q + 1 'step' 2*q + 1 'until' s 'do'
. . . . . . . l[r] := 0;
. . . . . . 'if' m '/' (2*q + 1) * (2*q + 1) < m 'and'
. . . . . . . n '/' (2*q + 1) * (2*q + 1) < n 'then' l[q] := 0
. . . . . . 'end';
. . . .
. . . 'if' m/'2*2 = m 'or' n/'2*2 = n 'then'
. . . . 'begin'
. . . . . a := 1;
. . . . . t[0,a] := 2;
. . . . . t[1,a] := ord(2,m);
. . . . . t[2,a] := ord(2,n)
. . . . 'end''else' a := 0;
. . . .
. . . 'for' q := 1 'step' 1 'until' s 'do'
. . . . 'if' l[q] = 1 'then'

```

```

. . . . 'begin'
. . . . . a := a + 1;
. . . . . t[0,a] := 2*q + 1;
. . . . . t[1,a] := ord(2*q + 1,m);
. . . . . t[2,a] := ord(2*q + 1,n)
. . . . 'end';
. . . .
. . . 'if' m > 1 'and' n = 1 'then'
. . . 'begin'
. . . . a := a + 1;
. . . . . t[0,a] := m;
. . . . . t[1,a] := 1;
. . . . . t[2,a] := 0
. . . 'end';
. . . .
. . . 'if' m = 1 'and' n > 1 'then'
. . . 'begin'
. . . . a := a + 1;
. . . . . t[0,a] := n;
. . . . . t[1,a] := 0;
. . . . . t[2,a] := 1
. . . 'end';
. . . .
. . . 'if' m > 1 'and' n > m 'then'
. . . 'begin'
. . . . a := a + 2;
. . . . . t[0,a-1] := m; t[0,a] := n;
. . . . . t[1,a-1] := 1; t[1,a] := 0;
. . . . . t[2,a-1] := 0; t[2,a] := 1
. . . 'end';
. . . .
. . . 'if' n > 1 'and' m > n 'then'
. . . 'begin'
. . . . a := a + 2;
. . . . . t[0,a-1] := n; t[0,a] := m;
. . . . . t[1,a-1] := 0; t[1,a] := 1;
. . . . . t[2,a-1] := 1; t[2,a] := 0
. . . 'end';
. . . .
. . . 'if' m > 1 'and' m = n 'then'
. . . 'begin'
. . . . a := a + 1;
. . . . . t[0,a] := m;
. . . . . t[1,a] := 1;
. . . . . t[2,a] := 1
. . . 'end';
. . 'end';
. 'end';
.
. output(4,<<<<gib den modus ein;>>>>);
. output(4,<<<<0 bei dateneingabe ueber eine datei auf den kanal 1>>>>);
. output(4,<<<<1 bei eingabe einzelner gleichungen>>>>);
. output(4,<<<<2 bei eingabe aller gleichungen ueber konsole>>>>);
. read(v); weiter:
. 'if' v = 0 'then' input(1,<<zz0>>,e);
. 'if' v = 1 'then' e := 1;
. 'if' v = 2 'then''begin'

```

```

. . output(4,<</,<<Wieviele Gleichungen werden eingegeben ?>>>>);
. . read(e) 'end';
. .
. 'begin'
. . 'integer''array' a[1:5,1:e]; 'integer' i,k,m,n,jz,jn;
. . 'boolean' minimal;
. .
. . 'if' v > 0 'then'
. . 'begin'
. . . output(4,<</,<<gib die Koeffizienten so ein>>>>);
. . . output(4,<</,<<a1 a3 a2 a4 a6>>>>);
. . . 'for' k := 1 'step' 1 'until' e 'do'
. . .   'for' i := 1, 2, 3, 4, 5 'do' read(a[i,k]);
. . 'end''else'
. . 'begin'
. . . 'for' k := 1 'step' 1 'until' e 'do'
. . .   input(1,<<d,d,-d,-3zd,-4zd>>,a[1,k],a[2,k],a[3,k],a[4,k],a[5,k]);
. . 'end';
. .
. . k := 0;
mark: k := k + 1;
. .
. . a1 := w[1] := a[1,k];
. . a3 := w[2] := a[2,k];
. . a2 := w[3] := a[3,k];
. . a4 := w[4] := a[4,k];
. . a6 := w[5] := a[5,k];
. . b2 := a1*a1 + 4*a2;
. . b4 := a1*a3 + 2*a4;
. . b6 := a3*a3 + 4*a6;
. . b8 := b2*a6 - a1*a3*a4 + a2*a3*a3 - a4*a4;
. . c4 := b2*b2 - 24*b4;
. . c6 := -b2*b2*b2 + 36*b2*b4 - 216*b6;
. . d := -b2*b2*b8 - 8*b4*b4*b4 - 27*b6*b6 + 9*b2*b4*b6;
. .
. . output(4,<<//,<<=====>>>>);
. .
. . output(4,<</,<<YY >>>>);
. . 'if' a1 = 0 'then''goto' m3;
. . sig(a1); output(abs(a1)); output(4,<<<<*XY >>>>);
m3: 'if' a3 = 0 'then''goto' m2;
. . sig(a3); output(abs(a3)); output(4,<<<<*Y >>>>);
m2: output(4,<<<<= XXX >>>>);
. . 'if' a2 = 0 'then''goto' m4;
. . sig(a2); output(abs(a2)); output(4,<<<<*XX >>>>);
m4: 'if' a4 = 0 'then''goto' m6;
. . sig(a4); output(abs(a4)); output(4,<<<<*X >>>>);
m6: 'if' a6 = 0 'then''goto' disk;
. . sig(a6); output(abs(a6));
. .
disk: 'if' d = 0 'then'
. . 'begin' output(4,<</,<<s i n g u l a e r !>>>>);
. . . 'if' c4 = 0 'then' output(4,<<<< ( Spitze )>>>>)
. . .   'else' output(4,<<<< ( Knoten )>>>>);
. . . 'goto' stop
. . 'end';
. .

```

```

. . m := max(1,abs(c4)); n := abs(d);
. .
. . tp(m,n,b,t);
. .
. . output(4,<<//,<<J - Invariante:>>>>);
. . output(4,<</,<<----->>>>);
. . output(4,<</,<<j = >>>>); sig(c4/d);
. . 'if' c4 = 0 'then'
. . . 'begin' output(4,<<<<0 >>>>); 'goto' delta 'end';
. . jz := jn := 1;
. . 'for' i := 1 'step' 1 'until' b 'do'
. . 'begin'
. . . 'integer' exp;
. . . exp := 3*t[1,i] - t[2,i];
. . . 'if' exp = 0 'then' goto inf;
. . . out(t[0,i]);
. . . output(4,<<<<*[>>>>);
. . . sig(exp); out(abs(exp));
. . . output(4,<<<<[ * >>>>);
. . . 'if' exp > 0 'then' jz := jz*t[0,i]**exp
. . . . 'else' jn := jn*t[0,i]**(-exp);
. . inf:
. . 'end';
. . output(4,<<<<= >>>>); sig(c4/d);
. . out(jz); output(4,<<<</>>>>); out(jn);
. .
delta: output(4,<<//,<<Diskriminante:>>>>);
. . output(4,<</,<<----->>>>);
. . output(4,<</,<<Delta = >>>>); sig(d);
. . 'for' i := 1 'step' 1 'until' b 'do'
. . 'if' t[2,i] > 0 'then'
. . 'begin'
. . . out(t[0,i]);
. . . output(4,<<<<*[>>>>);
. . . out(t[2,i]);
. . . output(4,<<<<[ * >>>>);
. . 'end';
. . output(4,<<<<= >>>>);
. . sig(d); out(abs(d));
. .
. . outout(4,<</,<<singulaere Fasern ueber :>>>>);
. . 'for' i := 1 'step' 1 'until' b 'do'
. . . 'if' t[2,i] > 0 'then'
. . . . 'begin'
. . . . . output(4,<</>>); out(t[0,i]);
. . . . . 'if' c4 = 0 'then'
. . . . . . output(4,<<<< (+) instabile faser, potentiell gut>>>>)
. . . . . . 'else'
. . . . . . 'begin'
. . . . . . . 'if' t[1,i] = 0 'then'
. . . . . . . . output(4,<<<< (*) semistabile Faser>>>>);
. . . . . . . 'if' 3*t[1,i] >= t[2,i] 'then'
. . . . . . . . output(4,<<<< (+) instabile Faser, potentiell gut>>>>);
. . . . . . . 'if' t[1,i] > 0 'and' 3*t[1,i] < t[2,i] 'then'
. . . . . . . . output(4,<<<< (+*) instabile Faser, potentiell schlecht>>>>);
. . . . . . . 'end';
. . . . . 'end';
. . . 'end';

```

```

. . .
. . minimal := 'true'; 'comment'
. . 'for' i := 1 'step' 1 'until' b 'do'
. . . 'if' t[2,i] > 0 'then'
. . . 'begin'
. . . . 'if' t[0,i] > 3 'and' t[1,i] > 3 'and' t[2,i] > 11 'then'
. . . . 'if' c4 = 0 'then'
. . . . 'begin'
. . . . . 'if' t[0,i] > 3 'and' t[2,i] > 11 'then'
. . . . . minimal := 'false'; 'comment'
. . . . . 'if' t[0,i] = 3 'and' t[2,i] > 11 'then'
. . . . .
. . . . .
. . . . 'if' t[0,i] > 3 'and' t[1,i] > 3 'and' t[2,i] > 11 'then'
. . . . minimal := 'false';
stop: . 'if' k < e 'then''goto' mark;
. 'end'
. output(4,<<//,<<=====>>>>);
. 'if' v > 0 'then'
. 'begin'
. . output(4,<<//,<<Sollen weitere Gleichungen gerechnet werden ?>>>>);
. . output(4,<<<< (nein = 0, ja = 1)>>>>);
. . read(v);
. . 'if' v = -1 'then''begin' v := 0; 'goto' weiter 'end';
. . 'if' v > 0 'then''goto' weiter;
. 'end';
'end' Version vom 11.3.80

```

4.3. Elliptic Code in C. The previous Algol program has been ported into C (still without the reduction routine for *non-minimal* WEIERSTRASS equations, which is incomplete in the Algol program). The only difference is that in Algol it is possible to nest procedures, which allows dynamical definition¹ of arrays depending on parameters. The procedure `tp(m,n,a,t)` for the prime factorisation in the Algol program does this quite clever, by decomposing c_4 and Δ in the same pass, which makes it easy to write the j -invariant in reduced form (canceling common factors).

The corresponding procedure here is `fact(N, d[], p[], e[])`, where N is factorised into the vector `p[]` of primes with exponents in the vector `e[]`: $N = p_1^{e_1} \cdots p_t^{e_t}$.

The array `d[]` consists of the test primes with $d[0] = 2, d[1] = 3, \dots, d[155612] = 2097169$, exhausting all primes $p < 2^{21}$ (we have in fact that $p = 2097169 > 2^{21}$). The routine to generate the prime divisors uses the sieve of ERATOSTHENES (see section 2), the factorisation into primes uses an algorithm of KNUTH (*The Art of Computer Programming* [1, vol. 2, §4.5.4], Algorithm A).

/*

/* -----

Module: ell.c

Description:

The Tate formulas for plane Weierstrass models

¹in the advent of C99 this is now also feasible in C - it remains to be added in my program

of elliptic curves over rational integers

Input: coefficients a1, a3, a2, a4, a6 (rational integers)

Output: prime factor decomposition of c4, c6, D and j
 reduction fiber type at bad reduction
 (Neron fiber type -- to be added) (conductor -- to be added),

Subroutines: sig, sign, display, sieve, fact, fiber

\$Date: 2012-12-02 16:28:11 +0100 (Sun, 02 Dec 2012) \$

\$Revision: 316 \$

Version: 1.7

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 distributed under the License is distributed on an "AS IS" BASIS,
 WITHOUT WARRANTIES OR CONDITIONS OF ANY KIND, either express or implied.
 See the License for the specific language governing permissions and
 limitations under the License.

-----*/

```
#include <stdio.h>
#include <stdlib.h>
#include <inttypes.h>

int main(void)
{
// constants

#define L      1024           // 2^10
#define M      1048584       // 2^20 + 8
#define YES    1             // flag for primes
#define NO     0             // flag for composite numbers

// variables

int s,t,i;
long a1, a3, a2, a4, a6;
long long b2, b4, b6, b8, c4, c6, d6, d8, D;
unsigned long p[30];        // array of prime factors
unsigned long q[30];        // array of bad primes
int e[30];                 // array of exponents
int f[30];                 // array of exponents
unsigned long d[155620];    // array of test divisors

// prototype statements for functions
```

```

int sig(long long);
int sign(long long);
int sieve(unsigned long []);
int fact(unsigned long long, unsigned long [], unsigned long [], int []);
int fiber(unsigned long, int, unsigned long [], int []);
void display(long long, unsigned long[]);

// start of main processing

t = sieve(d); // initialize test divisors
// t = 155612, p = 2097169 > 2^21

printf("Enter the coefficients of the Weierstrass equation \n");
printf("yy + a_1.xy + a_3.y = xxx + a_2.xx + a_4.x + a_6 \n\n");

scanf("%d %d %d %d %d", &a1, &a3, &a2, &a4, &a6);

// the Weierstrass-Tate formulaire

b2 = a1*a1 + 4*a2;
b4 = a1*a3 + 2*a4;
b6 = a3*a3 + 4*a6;
b8 = b2*a6 - a1*a3*a4 + a2*a3*a3 - a4*a4;

c4 = b2*b2 - 24*b4;
c6 = - b2*b2*b2 + 36*b2*b4 - 216*b6;

d6 = b2*b4 - 18*b6;
d8 = b2*b2*b4 - 48*b4*b4 + 18*b2*b6;

D = - b2*b2*b8 + b4*b4*b4 + 36*b4*b8 - 27*b6*b6 ;

// display the Weierstrass equation of the curve

printf("\nyy ");
if (a1) {sig(a1); printf(" %dxy ",abs(a1));}
if (a3) {sig(a3); printf(" %dy ",abs(a3));}
printf("= xxx ");
if (a2) {sig(a2); printf(" %dxx ",abs(a2));}
if (a4) {sig(a4); printf(" %dx ",abs(a4));}
if (a6) {sig(a6); printf(" %d ",abs(a6));}

if (D) {
printf(" elliptic curve\n");
// printf("j - invariant: %d / %d \n", c4*c4*c4,D);

printf("\nc4 = ");
display(c4,d);

printf("\nc6 = ");
display(c6,d);

printf("\nD = ");
display(D,d);

// fiber type at bad reduction primes

```

```

printf("\n\n** reduction fiber type ** \n");

t = fact(sign(c4)*c4, d, p, e);
t = fact(sign(D)*D, d, q, f);

// in q[1], ... , q[t] are the different primes of bad reduction

for (t=1; t<=q[0]; t++)
  switch (fiber(q[t],f[t],p,e)) {
    case +1:
      printf("unstable fiber at %u (cusp), potentially good \n", q[t]);
      break;
    case -1:
      printf("unstable fiber at %u (cusp), potentially bad", q[t]);
      printf(" ... and more details here\n");
      break;
    case 0:
      printf("semi-stable fiber at %d (node)\n", q[t]);
      break;
  }
}
else {
  printf(" singular curve (");
  if (c4)
    printf("node)\n");
  else
    printf("cusp)\n");
}

return 0;
}

/* -----
**
** Subroutines  sig, sign
**
** Description: sign of an integer (symbol resp. value)
**
**      Input:  N (integer)
**
**      Output: '+' or '-' resp. return value +1 or -1
**
** -----*/

int sig(long long N)
{
  if (N<0)
    printf("- ");
  else
    printf("+ ");
}

int sign(long long N)
{
  if (N<0)
    return -1;
}

```

```

    else
        return +1;
}

/* -----
**
** Subroutines  display
**
** Description:
**
**     Input:  x (integer), d (array of test divisors)
**
**     Output: prime factor decomposition in exponential notation
**
** -----*/

void display(long long x, unsigned long d[])
{
    int s,t;
    unsigned long q[30];
    int f[30];
    if (x) {
        t = fact(sign(x)*x, d, q, f);    // factorize the number
        printf("%PRId64" = ",x);
        sig(x);
        if (t) {
            for (s=1; s<t ; s++)
                printf("%u(%u) * ",q[s], f[s]);
            printf("%u(%u)",q[t], f[t]);
        }
        else
            printf("1");                // if no factors, it is 1
    } else
        printf("0");
}

/* -----
**
** Subroutine  sieve
**
** Description: int sieve(unsigned long d[]);
**
**     Generating array of test primes via sieve of Eratosthenes
**
**     Input:  none
**
**     Output: return value, vector d[] of test divisors
**
** -----*/

int sieve(unsigned long d[])
{
    // sieve of Eratosthenes

    unsigned int a, i;                // loop ind, exponent

```

```

/* the array prime[x] test for primality of 2*x + 1 */
char prime[M+1];          // byte array for flags

for (a = 0; a <= M; a++)  // initialize ...
    prime[a] = YES;      // ... to everything prime
for (a = 1; a <= L; a++)  // go thru all primes and
    if (prime[a])        // throw out the multiples
        for (i = 2*a*(a+1); i <= M; i += 2*a+1)
            prime[i] = NO; // multiples are not prime

// populate the vector of test divisors

i = 0;
d[i++] = 2;              // set the even prime

for (a = 1; a <= M; a++)
    if (prime[a])
        d[i++] = 2*a + 1; // set the uneven primes

return i;
}

/* -----
**
** Subroutine fact
**
** Description:
** The number N is factored into primes
**  $N = p[1]**e[1] * \dots * p[t]**e[t]$ 
**
** Input: N (integer), test divisors in vector d[]
**
** Output: vector p[], vector e[]
** the vector p[] contains the different primes
** the vector e[] contains the orders of N at these primes
** p[0] = t = return value = number of primes (also t = 0)
** (if N = 0 then e[0] = 1, otherwise e[0] = 0)
** -----*/

int fact(unsigned long long N, unsigned long d[], unsigned long p[], int e[])
{
// this is Knuth's algorithm, modified with exponents

unsigned int i, t, k;
unsigned long long n, q, r;

n = N; // A1 initialize
i = t = k = 0;

while (n > 1) { // A2 n = 1 ?
    q = n / d[k]; // A3 divide
    r = n - q * d[k];
    if (r == 0) { // A4 zero remainder ?
        if (i == 0)

```

```

        p[++] = d[k];           // A5 new factor found
        e[t] = ++i;           // increase exponent
        n = q;
    }                           // return to A2
    else {
        if (q > d[k]) {       // A6 low quotient ?
            k = k + 1;       // no, .. try next
            i = 0;
        }                       // return to A3
        else {                // yes, ..
            p[++] = n;       // A7 n is prime
            e[t] = 1;       // exponent is 1
            n = 1;         // terminate
        }
    }
}
p[0] = t;                       // number of factors
if (n)
    e[0] = 0;                   // e[0] discriminates
else
    e[0] = 1;                   // .. if n vanishes

return t;
}

/* -----
**
** Subroutine  fiber
**
** Description:
**     fiber returns the fiber type
**         +0     semi-stable fiber (node)
**         +1     unstable fiber (cusp), but potentially good
**         -1     unstable fiber (cusp), and potentially bad
**
** Input:  prime q (of bad reduction), exponent f (of D)
**         vector p[], exponent vector e[] (of c4)
**
** Output: return value is fiber type
**
** -----*/

int fiber(unsigned long q, int f, unsigned long p[], int e[])
{
    unsigned int t;
    int type = 0;                // default type ss

    if (e[0])
        type = 1;               // if c4 = 0
    else {
        for (t=1;t<=p[0];t++)    // test for common prime
            if (q == p[t]) {    // if found,
                type = sign(3*e[t]-f); // set type accordingly
                break;
            }
    }
}
return type;

```

```

}
/*
*/

```

4.4. Calculation of a resultant. In fact, in this section there is no computer algorithm, only a normal calculation of a resultant.

I make use of TATE's *formulaire*, as it is used by TATE in [8] or in the appendix to LANG [3], or in his letter to CASSELS in [2]. See also SILVERMAN [7, chap. III, §1].

In the formula for duplication of points on an elliptic curve we have for $P = (x, y)$ and $[2]P = (x_2, y_2)$:

$$x_2 = \frac{x^4 - b_4x^2 - 2b_6x - b_8}{4x^3 + b_2x^2 + 2b_4x + b_6}$$

(see SILVERMAN [7, chap. III, §2, 2.3d]).

We are going to calculate the resultant $R = R((1, 0, -b_4, -2b_6, -b_8), (4, b_2, 2b_4, b_6))$ of the two polynomials occurring in the numerator and denominator. We will simplify this determinant before multiplying it out, and also make use of $4b_8 = b_2b_6 - b_4^2$ at times.

The resultant is (cf. LANG *Algebra* [4, IV. §8]):

$$\begin{aligned}
R &= \begin{vmatrix} 1 & 0 & -b_4 & -2b_6 & -b_8 & 0 & 0 \\ 0 & 1 & 0 & -b_4 & -2b_6 & -b_8 & 0 \\ 0 & 0 & 1 & 0 & -b_4 & -2b_6 & -b_8 \\ 4 & b_2 & 2b_4 & b_6 & 0 & 0 & 0 \\ 0 & 4 & b_2 & 2b_4 & b_6 & 0 & 0 \\ 0 & 0 & 4 & b_2 & 2b_4 & b_6 & 0 \\ 0 & 0 & 0 & 4 & b_2 & 2b_4 & b_6 \end{vmatrix} = \\
&= \begin{vmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ -4 & -b_2 & -6b_4 & 1 & 0 & 0 & 0 \\ 0 & -4 & -b_2 & 0 & 1 & 0 & 0 \\ 0 & 0 & -4 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{vmatrix} \begin{vmatrix} 1 & 0 & -b_4 & -2b_6 & -b_8 & 0 & 0 \\ 0 & 1 & 0 & -b_4 & -2b_6 & -b_8 & 0 \\ 0 & 0 & 1 & 0 & -b_4 & -2b_6 & -b_8 \\ 4 & b_2 & 2b_4 & b_6 & 0 & 0 & 0 \\ 0 & 4 & b_2 & 2b_4 & b_6 & 0 & 0 \\ 0 & 0 & 4 & b_2 & 2b_4 & b_6 & 0 \\ 0 & 0 & 0 & 4 & b_2 & 2b_4 & b_6 \end{vmatrix} = \\
&= \begin{vmatrix} b_2b_4 + 9b_6 & 3b_2b_6 + 5b_4^2 & b_2b_8 + 12b_4b_6 & 6b_4b_8 \\ 6b_4 & b_2b_4 + 9b_6 & 3b_2b_6 - b_4^2 & b_2b_8 \\ b_2 & 6b_4 & 9b_6 & 4b_8 \\ 4 & b_2 & 2b_4 & b_6 \end{vmatrix}
\end{aligned}$$

We simplify this slightly and develop the determinant according to its first row:

$$\begin{aligned}
R &= \begin{vmatrix} 1 & 0 & -b_4 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{vmatrix} \begin{vmatrix} b_2b_4 + 9b_6 & 3b_2b_6 + 5b_4^2 & b_2b_8 + 12b_4b_6 & 6b_4b_8 \\ 6b_4 & b_2b_4 + 9b_6 & 3b_2b_6 - b_4^2 & b_2b_8 \\ b_2 & 6b_4 & 9b_6 & 4b_8 \\ 4 & b_2 & 2b_4 & b_6 \end{vmatrix} \\
&= \begin{vmatrix} 9b_6 & 3b_2b_6 - b_4^2 & b_2b_8 + 3b_4b_6 & 2b_4b_8 \\ 6b_4 & b_2b_4 + 9b_6 & 3b_2b_6 - b_4^2 & b_2b_8 \\ b_2 & 6b_4 & 9b_6 & 4b_8 \\ 4 & b_2 & 2b_4 & b_6 \end{vmatrix} = \\
&= 9b_6 \begin{vmatrix} b_2b_4 + 9b_6 & 3b_2b_6 - b_4^2 & b_2b_8 \\ 6b_4 & 9b_6 & 4b_8 \\ b_2 & 2b_4 & b_6 \end{vmatrix} - (3b_2b_6 - b_4^2) \begin{vmatrix} 6b_4 & 3b_2b_6 - b_4^2 & b_2b_8 \\ b_2 & 9b_6 & 4b_8 \\ 4 & 2b_4 & b_6 \end{vmatrix} + \\
&+ (b_2b_8 + 3b_4b_6) \begin{vmatrix} 6b_4 & b_2b_4 + 9b_6 & b_2b_8 \\ b_2 & 6b_4 & 4b_8 \\ 4 & b_2 & b_6 \end{vmatrix} - 2b_4b_8 \begin{vmatrix} 6b_4 & b_2b_4 + 9b_6 & 3b_2b_6 - b_4^2 \\ b_2 & 6b_4 & 9b_6 \\ 4 & b_2 & 2b_4 \end{vmatrix}
\end{aligned}$$

Writing out the determinants

$$\begin{aligned}
R &= 9b_6(9(b_2b_4 + 9b_6)b_6^2 + 4(3b_2b_6 - b_4^2)b_2b_8 + 12b_2b_4^2b_8 - 9b_2^2b_6b_8 - \\
&- 6(3b_2b_6 - b_4^2)b_4b_6 - 8(b_2b_4 + 9b_6)b_4b_8) - \\
&- (3b_2b_6 - b_4^2)(54b_4b_6^2 + 16(3b_2b_6 - b_4^2)b_8 + 2b_2^2b_4b_8 - \\
&- 36b_2b_6b_8 - (3b_2b_6 - b_4^2)b_2b_6 - 48b_4^2b_8) + \\
&+ (b_2b_8 + 3b_4b_6)(36b_4^2b_6 + 16(b_2b_4 + 9b_6)b_8 + b_2^3b_8 - \\
&- 24b_2b_4b_8 - (b_2b_4 + 9b_6)b_2b_6 - 24b_2b_4b_8) - \\
&- 2b_4b_8(72b_4^3 + 36(b_2b_4 + 9b_6)b_6 + b_2^2(3b_2b_6 - b_4^2) - \\
&- 24(3b_2b_6 - b_4^2)b_4 - 2b_2b_4(b_2b_4 + 9b_6) - 54b_2b_4b_6) = \\
&= 9b_6(3b_2^2b_6b_8 - 9b_2b_4b_6^2 + 6b_4^3b_6 - 72b_4b_6b_8 + 81b_6^3) - \\
&- (3b_2b_6 - b_4^2)(2b_2^2b_4b_8 - 3b_2^2b_6^2 + b_2b_4^2b_6 + 12b_2b_6b_8 - 64b_4^2b_8 + 54b_4b_6^2) + \\
&+ (b_2b_8 + 3b_4b_6)(b_2^3b_8 - b_2^2b_4b_6 - 32b_2b_4b_8 - 9b_2b_6^2 + 36b_4^2b_6 + 144b_6b_8) - \\
&- 2b_4b_8(3b_2^3b_6 - 3b_2^2b_4^2 - 108b_2b_4b_6 + 96b_4^3 + 324b_6^2) =
\end{aligned}$$

use $4b_8 = b_2b_6 - b_4^2$

$$\begin{aligned}
&= 9b_6(3b_2^2b_6b_8 - 27b_2b_4b_6^2 + 24b_4^3b_6 + 81b_6^3) - \\
&- (3b_2b_6 - b_4^2)(2b_2^2b_4b_8 - 18b_2b_4^2b_6 + 16b_4^4 + 54b_4b_6^2) + \\
&+ (b_2b_8 + 3b_4b_6)(b_2^3b_8 - 9b_2^2b_4b_6 + 8b_2b_4^3 + 27b_2b_6^2) - \\
&- 6b_2^3b_4b_6b_8 + 6b_2^2b_4^3b_8 + 54b_2^2b_4^2b_6^2 - 102b_2b_4^4b_6 - 162b_2b_4b_6^3 + 48b_4^6 + 162b_4^3b_6^2 = \\
&= 27b_2^2b_6^2b_8 - 243b_2b_4b_6^3 + 216b_4^3b_6^2 + 729b_6^4 - \\
&- 6b_2^3b_4b_6b_8 + 2b_2^2b_4^3b_8 + 54b_2^2b_4^2b_6^2 - 66b_2b_4^4b_6 - 162b_2b_4b_6^3 + 16b_4^6 + 54b_4^3b_6^2 + \\
&+ b_2^4b_8^2 - 6b_2^3b_4b_6b_8 - 27b_2^2b_4^2b_6^2 + 8b_2^2b_4^3b_8 + 24b_2b_4^4b_6 + 27b_2^2b_6^2b_8 + 81b_2b_4b_6^3 - \\
&- 6b_2^3b_4b_6b_8 + 6b_2^2b_4^3b_8 + 54b_2^2b_4^2b_6^2 - 102b_2b_4^4b_6 - 162b_2b_4b_6^3 + 48b_4^6 + 162b_4^3b_6^2 = \\
&= b_2^4b_8^2 - 18b_2^3b_4b_6b_8 + 81b_2^2b_4^2b_6^2 + 16b_2^2b_4^3b_8 + 54b_2^2b_6^2b_8 - \\
&- 144b_2b_4^4b_6 - 486b_2b_4b_6^3 + 432b_4^3b_6^2 + 64b_4^6 + 729b_6^4 = \Delta^2
\end{aligned}$$

where $\Delta = -b_2^2b_8 + 9b_2b_4b_6 - 8b_4^3 - 27b_6^2$ is the discriminant of the elliptic curve.

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